Eidgenössische Technische Hochschule Zürich
Swiss Federal Institute of Technology Zurich

Distributed Computing

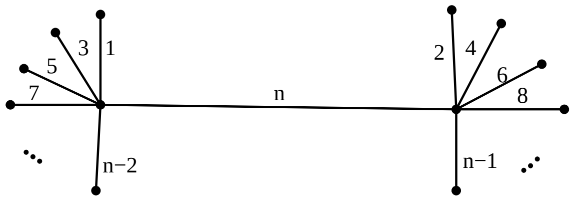
FS 2013

Prof. R. Wattenhofer
Stephan Holzer

Principles of Distributed Computing
Sample Solution to Exercise 7

1 Lightest Edges

a) Clearly, the execution of this algorithm cannot take more than n rounds. Let the $(n-1)$ lightest edges form two stars of the same size and the n^{th} lightest edge connect the two centers of the stars. We are not interested in the distribution of the other weights. In this scenario it takes 1 round until the $n-1$ lightest edges are announced. In the next $\lceil n/2 \rceil - 2$ rounds, the centers of the stars will announce these edges once more. Only in round $\lceil n/2 \rceil$ the two center nodes announce the n^{th} lightest edge. Since it is necessary to know this edge, the algorithm cannot terminate earlier and the time complexity of this algorithm is $\Omega(n)$.



b) We first prove that the time complexity is upper bounded by $\lceil \sqrt{2n} \rceil \in O(\sqrt{n})$. After $\lceil \sqrt{2n} \rceil$ rounds, all nodes with at most $\sqrt{2n}$ edges among the n lightest edges have broadcast all relevant edges known to them. That means, after $\lceil \sqrt{2n} \rceil$ rounds, there can only be missing edges between nodes that initially had at least $\sqrt{2n}+1$ lightest edges leading to nodes that are also incident to at least $\sqrt{2n}$ lightest edges. Assume there is such a node. Since each edge connects two nodes, initially we must have had at least $(\sqrt{2n}+1)\sqrt{2n}/2 > n$ lightest edges, a contradiction.

We now construct a worst-case example. Each edge connecting two nodes from a specific set of $\lfloor \sqrt{2n} \rfloor$ nodes is assigned one of the n smallest weights. Since there are $\binom{\lfloor \sqrt{2n} \rfloor}{2} \leq n$ edges between these nodes, we know that all edges between these nodes must be broadcast. Apparently, the $\lfloor \sqrt{2n} \rfloor$ nodes will announce at most the same number of edges in each round. Thus, in total at least $\lfloor \sqrt{2n} \rfloor / 2 \in \Omega(\sqrt{n})$ rounds are required.

c) Node v can send the n^{th} smallest edge weight to all nodes. Every node v_i can now determine how many among its edges (v_i, v_j), where $i < j$, belong to the n lightest edges and send this value N_i to all nodes. Now, the nodes know to which node they have to send their edge weights such that they can be distributed in the next round without contention: Node v_i sends its smallest weight to the node v_k, where $k = 1 + \sum_{j=1}^{i-1} N_j$, the next one to v_{k+1}, etc. Thus, every node receives exactly one edge weight to forward to all nodes. This procedure takes four rounds, i.e., the time complexity is $O(1)$.