


Theorem:  $n - m + f = 2$

Proof  
[Sketch]

$m=0$ :   $\Rightarrow n - m + f = 1 - 0 + 1 = 2 \checkmark$

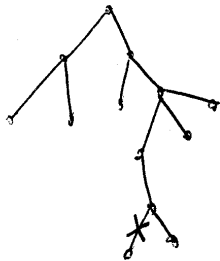
$m > 0$ : (assume formula correct for  $m-1$ )

Tree

remove leaf

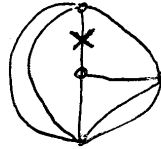
$$\Rightarrow n' = n - 1$$

$$m' = m - 1 \checkmark$$



Not tree

remove edge of cycle



$$\Rightarrow m' = m - 1$$

$$f' = f - 1 \checkmark$$

Theorem

Simple, connected, planar graph with  $n$  nodes  
has at most  $3n - 6$  edges ( $n \geq 3$ )

Proof

• each edge bounds at most 2 faces  
• each face bounded by at least 3 edges  $\Rightarrow 3f \leq 2m$

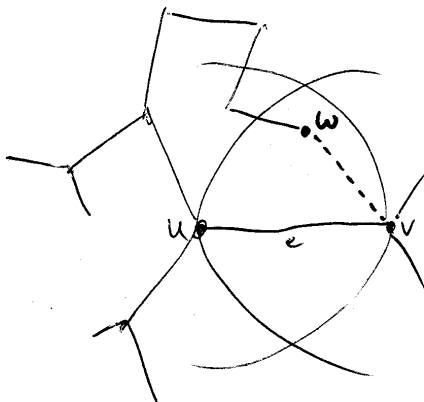
$$n - m + f = 2$$

$$3n - 3m + 3f = 6 \Leftrightarrow \underline{3n - 6} = 3m - 3f \geq 3m - 2m = m$$

MST  $\in$  RNG

Assume Contradiction:  $e \in \text{MST}$

$e \notin \text{RNG} \Rightarrow$  there is a point  $w$  in the interior (strictly)

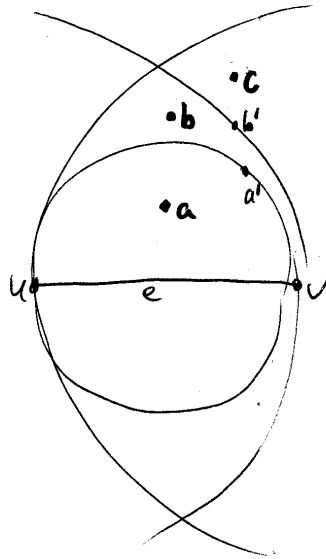


Remove  $e$  from MST  
 $\Rightarrow$  Two trees  $T_u, T_v$   
 $w$  is  $\in$  of  $\{T_u, T_v\}$   
 w.l.o.g  $w \in T_u$

We can reconnect  $T_u$  with  
 $T_v$  with the edge  $(v, w)$   
 better MST!  $\nabla$   $\blacksquare$

$$\underline{RNG} \subseteq \underline{GC}$$

by Definition



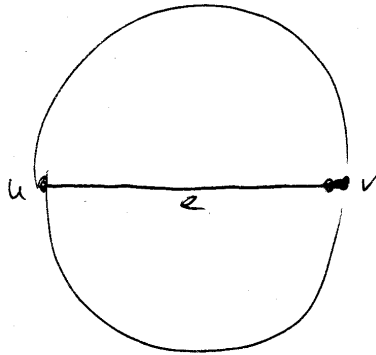
	<u>RNG</u>	<u>GC</u>
a)	$\neq$	$\neq$
b)	$\neq$	$\in$
c)	$\in$	$\in$

Comment:

a')	$\neq$	$\neq$
b')	$\in$	$\in$

GG  $\subseteq$  DT

by Definition

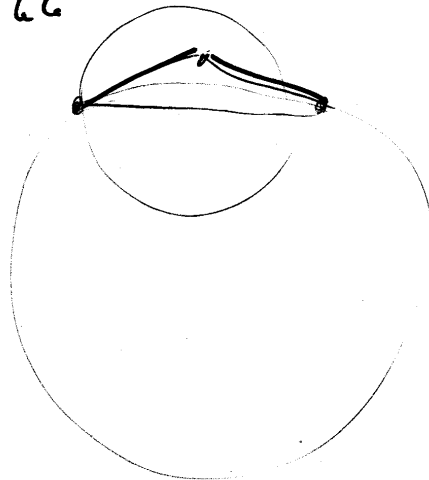


$e$  in GG if  
disk( $u, v$ ) contains no  
other node

DT:  $e$  in DT if  
any disk with  $u, v$   
on boundary contains  
no other node

Example:

DT, GG



MST connected

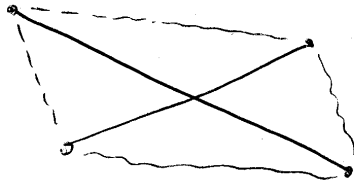
by definition

DT planar

by definition ... however, not quite so easy.

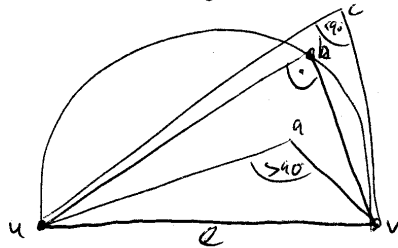
CG planar

Assume not



$= 360^\circ$

there is an angle  $\geq 90^\circ$



edge  $e$  exists  $\Rightarrow$  angle  $< 90^\circ$   $\hookrightarrow$   $\blacktriangleright$

GG contains Minimum Energy Path

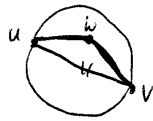
Proof:

Let this be MEP  
↓



Assume two nodes are not neighbors in GG.

Then there is a node  $w$  in the circle by  $u, v$ .



If  $uw$  or  $wv$  are not neighbors then you do the same again (recursively)

Otherwise  $E(u, w) + E(w, v) = \overline{uw}^\alpha + \overline{wv}^\alpha \leq \overline{uv}^\alpha$  (for  $\alpha \geq 2$ )

GG  $\cap$  UDG

Def: UDG

$e \in E$  of UDG  $\Leftrightarrow |e| \leq 1$ .

GG  $\cap$  UDG contains Minimum Energy Path

Proof as above, except the first path is MEP in UDG