Chapter 6
PEER-TO-PEER COMPUTING

Distributed
Computing
Group

Computer Networks
Summer 2005

Overview

- What is Peer-to-Peer?
- Dictionary
  - Distributed Hashing
  - Search
  - Join & Leave
- Other systems
- Conclusion

"Peer-to-Peer" is…

- Software: Napster, Gnutella, Kazaa, …
- File "sharing"
- Legal issues, RIAA
- Direct data exchange between clients
- Best effort, no guarantees
- 80% of Web Traffic "P2P"

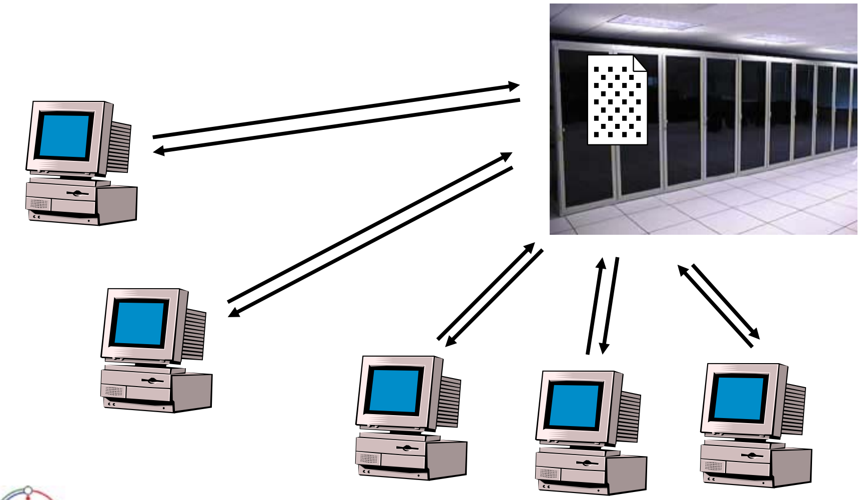
…a socio-cultural phenomenon!

"Peer-to-Peer" is also…

- A hot research area: Chord, Pastry, …
- A paradigm beyond Client/Server
- Dynamics (frequent joins and leaves)
- Fault tolerance
- Scalability
- Dictionary… and more!

… a new networking philosophy/technology!

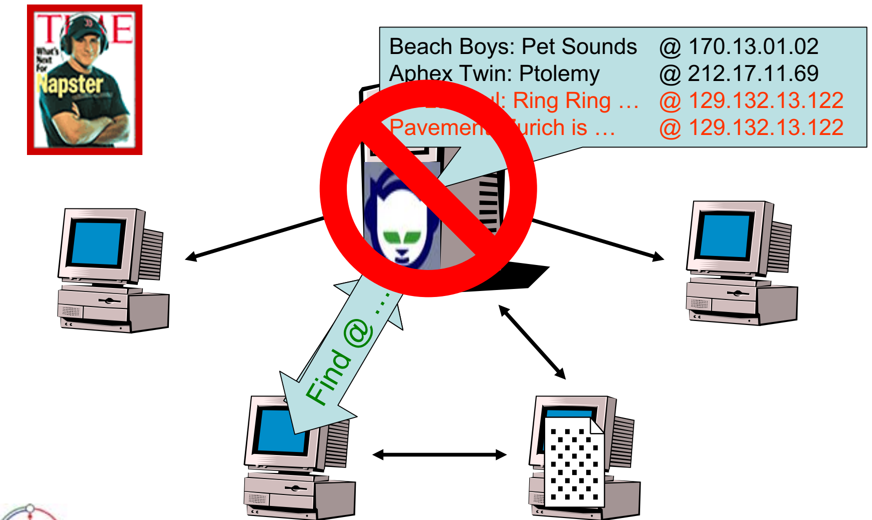
Client/Server



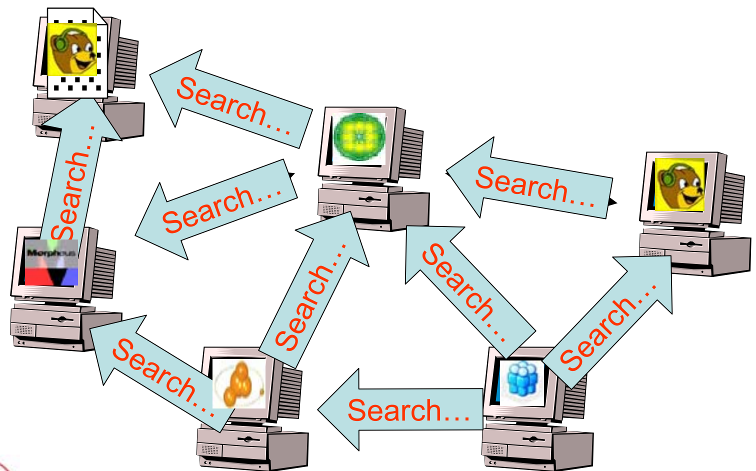
Client/Server Problems

- Scalability
  – Can server serve 100, 1’000, 10’000 clients?
  – What’s the cost?
- Security / Denial-of-Service
  – Servers attract hackers
- Replication
  – Replicating for security
  – Replicating close to clients (“caching”)

Case Study: Napster



Case Study: Gnutella



Pros/Cons Gnutella

- totally decentralized
- totally inefficient
  - “flooding” = directionless searching
- Gnutella often does not find searched item
  - TTL
  - Gnutella “not correct”

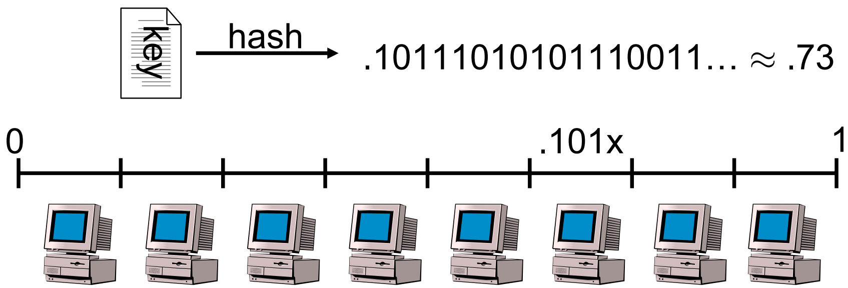
Dictionary ADT

- A collection of objects
  - Each object uniquely identified by key
- Supports these operations:
  - Search(key) → object(key)
  - Insert(key, object) → OK?
  - Delete(key) → OK?

Dictionary Implementations

- Classic Implementations
  - Search Tree (balanced, B-Tree)
  - Hashing (various forms)
- “Distributed” Implementations
  - Linear Hashing
  - Consistent Hashing

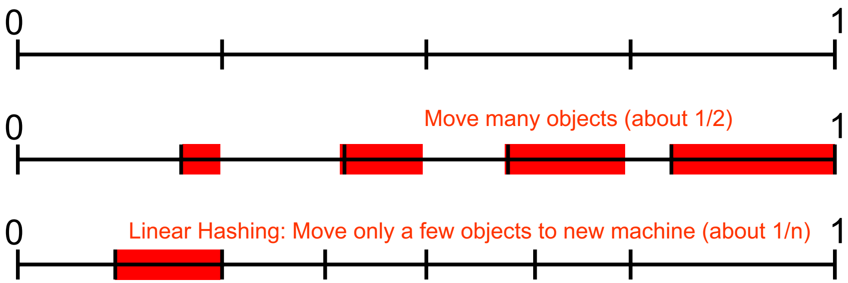
Distributed Hashing



- Remark: Instead of storing a document at the right peer, just store a forward-pointer

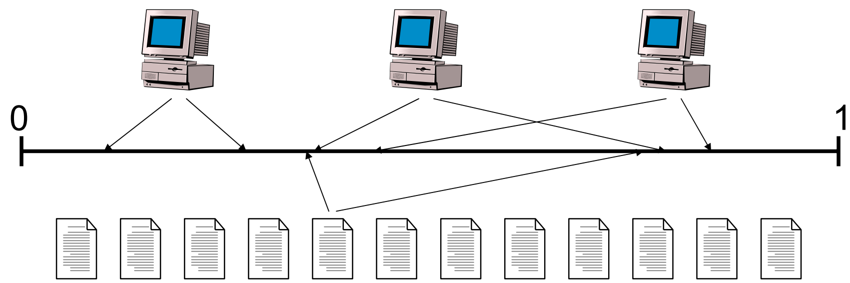
Linear Hashing

- Problem: More and more objects should be stored; need to buy new machines!
- Example: From 4 to 5 machines



Consistent Hashing

- Needs central dispatcher
- Idea: Also the machines get hashed! Each machine is responsible for the files closest to it. Use multiple hash funct. for reliability.



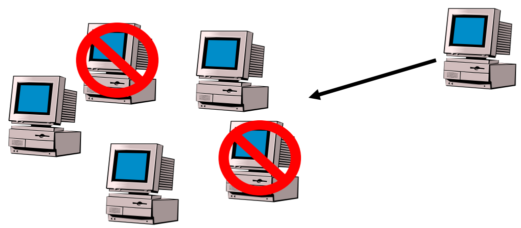
Not quite happy yet…

- Problem with both linear and consistent hashing is that all the participants of the system must know all peers…

- Number one challenge: Dynamics!
  - Peers join and leave

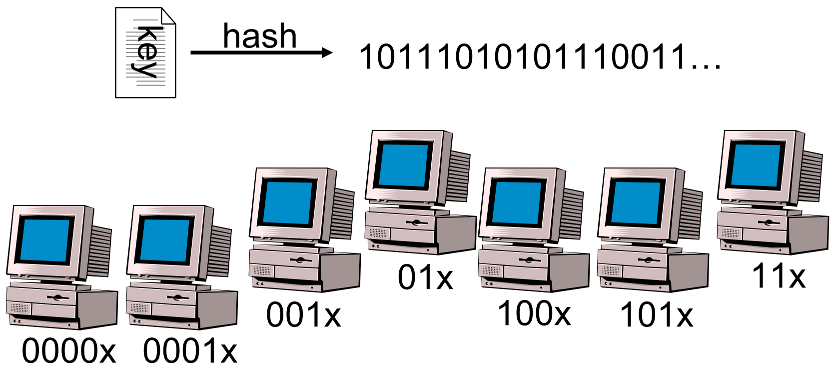
Dynamics

- Machines (peers) are unreliable
  - Joins; worse: spontaneous leaves!



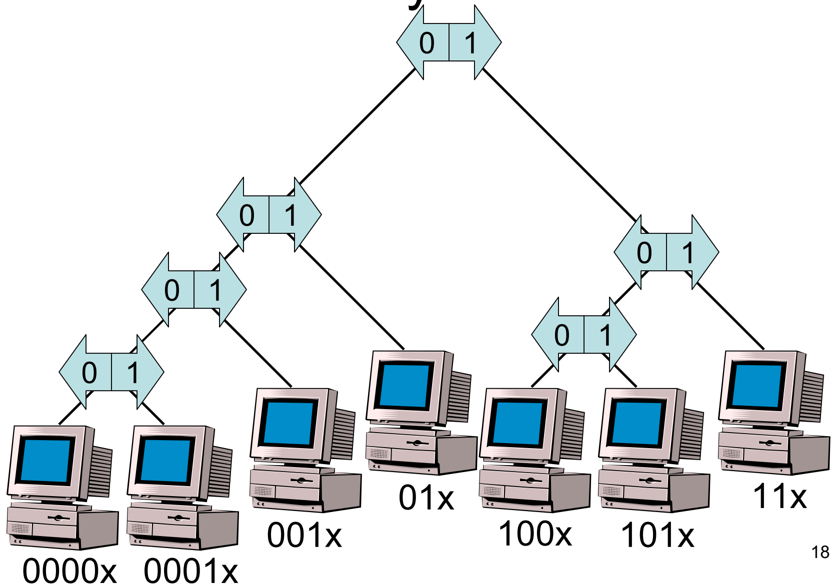
- Decentralized (“symmetric”) System
  - scalable, fault tolerant, dynamic

P2P Dictionary = Hashing



- Remark: Instead of storing a document at the right peer, just store a forward-pointer

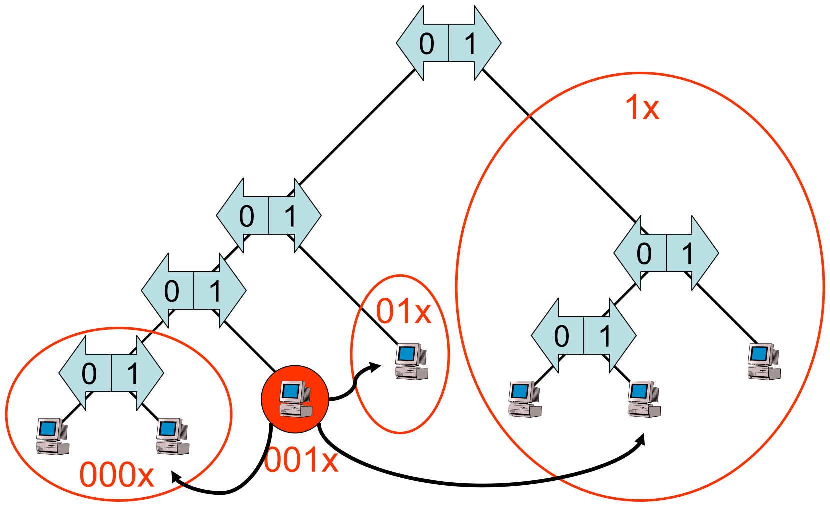
P2P Dictionary = Search Tree



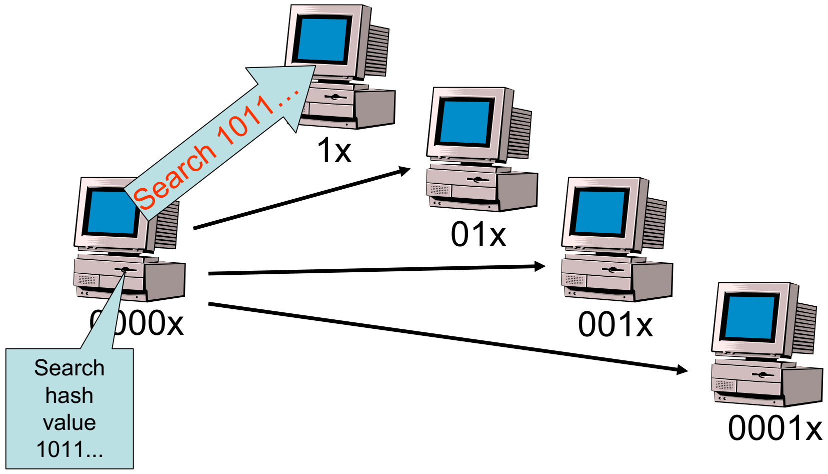
But who stores search tree?

- In particular, where is the root stored?
  - Root is scalability & fault tolerance problem
  - There is no root…!
- If a peer wants to store/search, how does it know where to go?
  - Does every peer know all others?
  - Dynamics! If a peer leaves, all peers must be notified. Too much overhead
  - Idea: Every peer only knows subset of others

The Neighbors of Peer 001x



P2P Dictionary: Search



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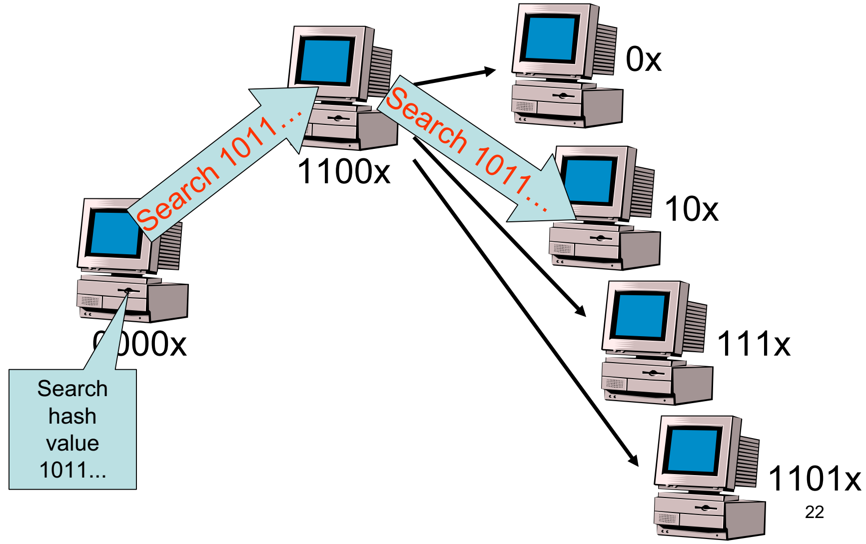
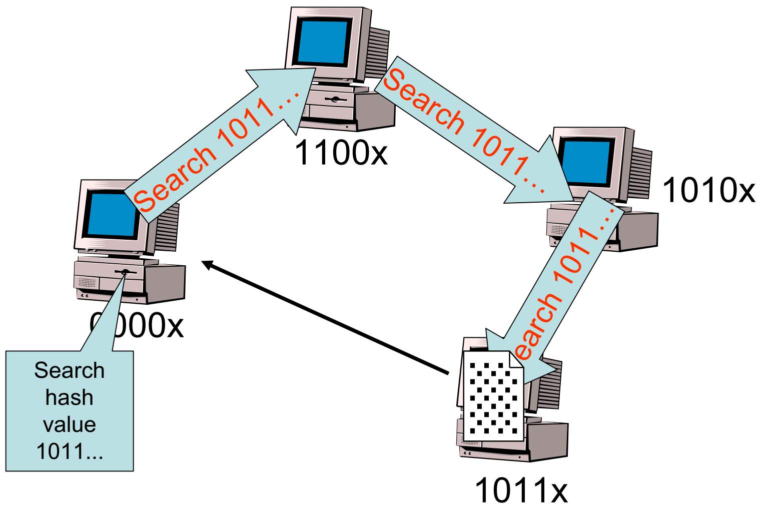
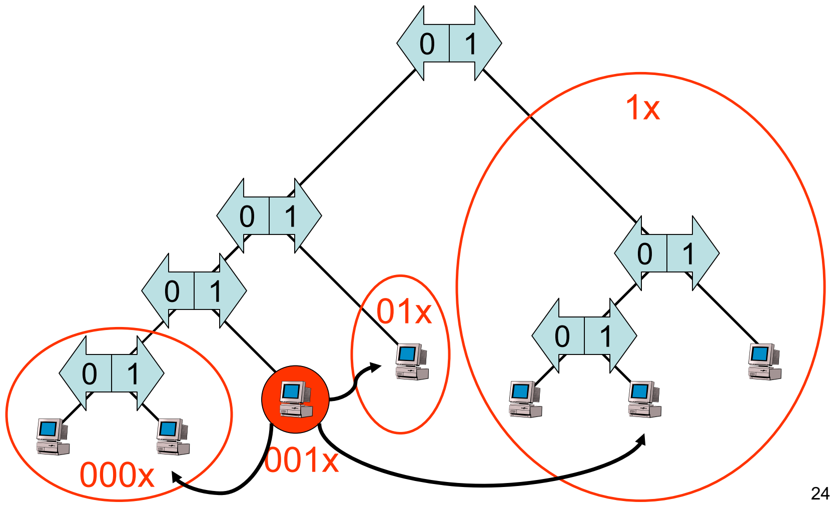


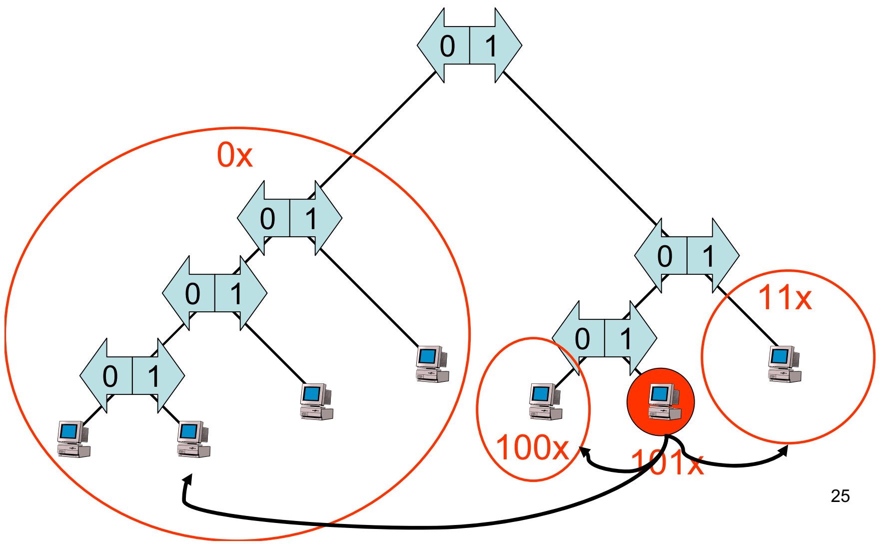
P2P Dictionary: Search



Again: 001 searches 100



001 searches 100 (continued)



Search Analysis

- We have n peers in system
- Assume that “tree” is roughly balanced
  – Leaves (peers) on level $\log_2 n \pm$ constant
- Search has O(log n) steps
  – After k’th step, you are in subtree on level k
  – A “step” is a UDP (or TCP) message
  – Latency is dependent on P2P size (world!)

Peer Join

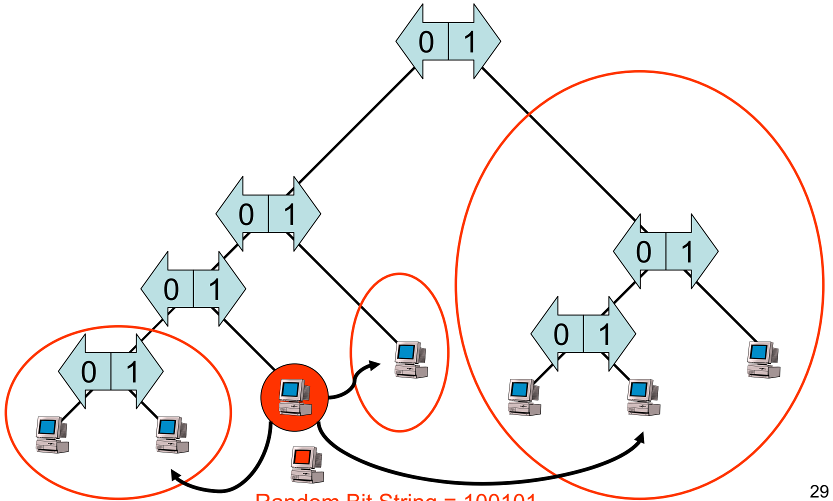
- Part 1: Bootstrap
- In order to join a P2P system, a joiner must already know a peer already in system. Typical solutions are
  – Ask a central authority for a list of IP addresses that have been in the P2P regularly; look up a listing on a web site
  – Try some of those you met last time
  – Just ping randomly (in the LAN)
- Part 2: Find your place in P2P system

2. Find your place

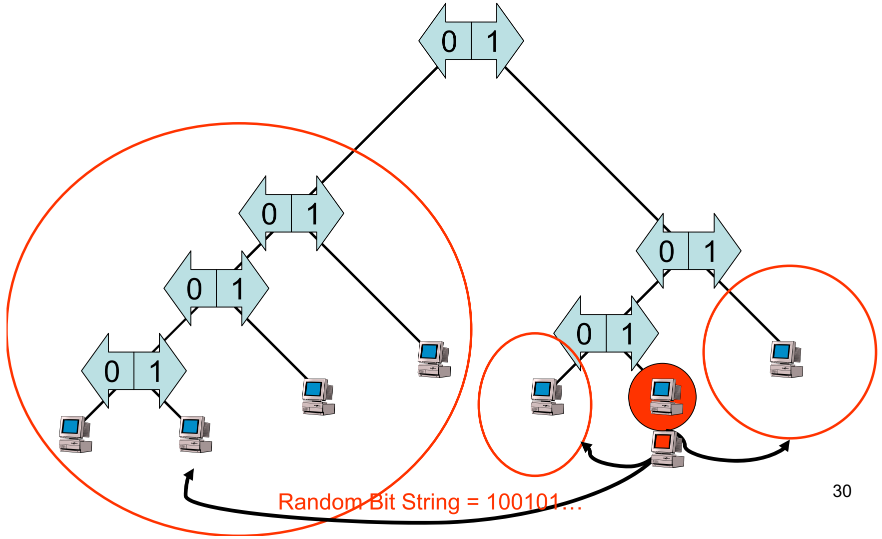
- The random method: Choose a random bit string (which determines the place)
- Search* for the bit string
- Split with the current leave responsible for the bit string
- Search* for your neighbors

* These are standard searches

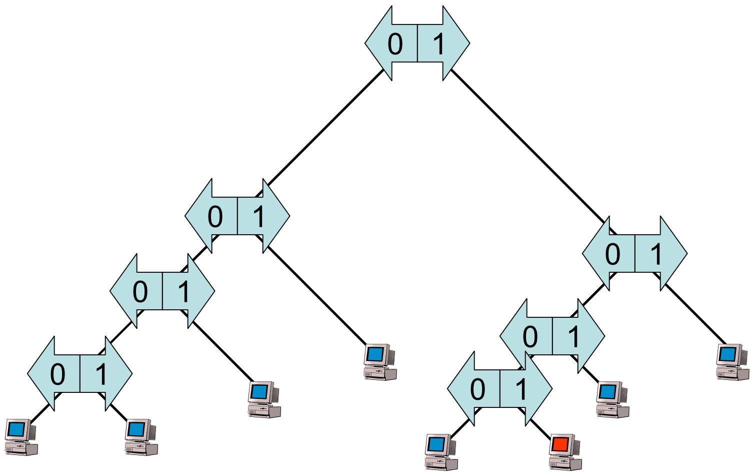
Example: Bootstrap with 001 peer



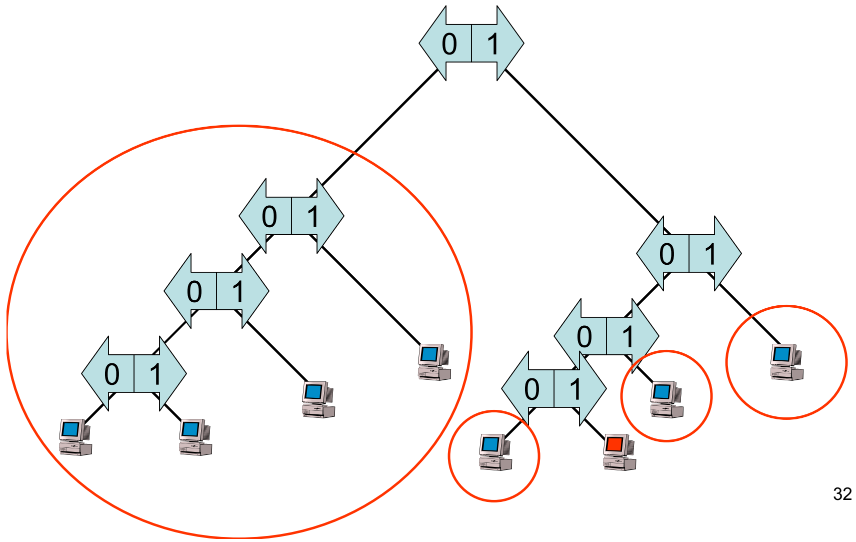
Joiner searches 100101…



Joiner found 100 leave → split



Find neighbors



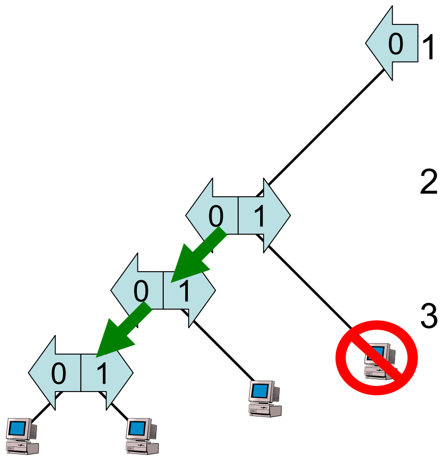
Random Join Discussion

- If tree is balanced, the time to join is
  - O(log n) for the first part
  - O(log n)·O(log n) = O($\log^2$ n) for the second part
- It is believe that since all the peers are chosen their position randomly, the tree will more or less be balanced.
  - However, theory and simulations show that this is widely believed but not really true.

Leave

- Since a leave might be spontaneous, it must be detected first. Naturally this is done by the neighbors in the P2P system (all peers periodically ping neighbors).
- If a peer that left was detected, it must be replaced. If peer had sibling leaf, the sibling might just do a "reverse split."
- If not, search recursively… example!

Peer 01 leaves spontaneously



1. Go down sibling tree, until you hit sibling leaves.
2. Make the left sibling the new common node.
3. Move the free right sibling to the empty spot.

Was that all?

- Yes, you now mastered all the P2P basics… Congratulations!
- But there are some nasty "technicalities" ☺
- Most importantly we would like to know what happened to the data that was stored at the peer that left (important question if we want to use the P2P network as a storage/file system). We study that soon…
- First some other comments…

Questions of experts…

- Q: I know so many other structured peer-to-peer systems; they are completely different from the one you showed us!

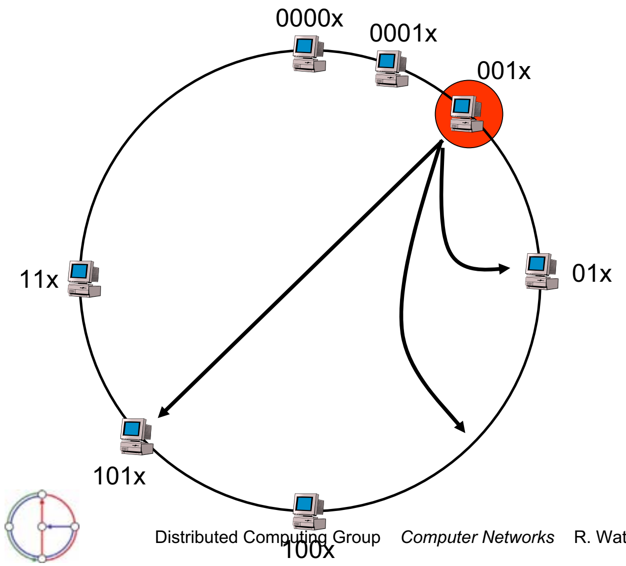
- A: They *look* different, but in fact the difference comes mostly from the way they are presented. (I give a few examples on the next slides)

Chord

- The most cited system by Ion Stoica, Robert Morris, David Karger, M. Frans Kaashoek, and Hari Balakrishnan, MIT, presented at ACM SIGCOMM 2001.
- Most discussed system in distributed systems and networking books, for example in Edition 4 of Tanenbaum’s Computer Networks.

- There are extensions on top of it, such as CFS, Ivy

Chord



- Every peer has log n many neighbors; one in about distance 2^{-k}, k=1, 2, …, log n

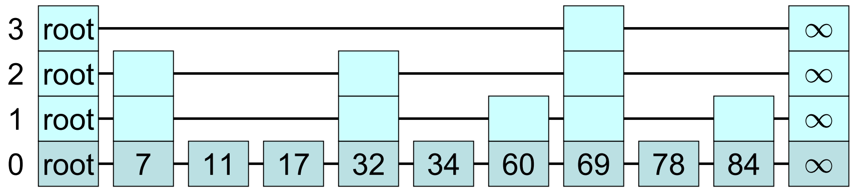
Skip List

- Are you afraid of programming balanced search trees (e.g. AVL or red-black tree)?!?

- Then the skip list is a data structure for you!

- Idea: Ordered linked list with extra pointers

Skip List



- (Doubly) linked list, with sorted items
- All items have additional pointers on levels 1, …, k, with probability 2^{-k}
- Search, insert, delete: Start with root, search for the right interval on highest level, then continue with lower levels.

Skip List

- It can easily be shown that search, insert, and delete terminate in O(log n) expected time, if there are n items in the skip list
- Also, on expectation, the number of pointers is only twice as many as with a regular linked list, thus the memory overhead is edible
- As a plus, the items are always ordered…

Skip Net

- Use the skip list as a peer-to-peer architecture: Again each peer gets a random value between 0 and 1, and is then responsible for storing that interval.
- Instead of a root and a sentinel node ("∞"), the list is short-wired as a ring
- There exist several proposals towards this end…

Many many others...

- Original work by Plaxton, Rajaraman, and Richa; "unfortunately" theory paper, so it includes many other aspects, such as a distance discussion… similar proposals are Pastry/Tapestry, or Kademlia.
- Some proposals improve the design; e.g. The Viceroy resp. Koorde proposals are Butterfly-based resp. DeBruijn-based and therefore only need a constant number of neighbors per peer.
- Closest/best design in reality is Freenet. However, Freenet has some questionable design properties

Why should I care?

- Q: I don't want to program a worldwide music stealing application, so why should I care?

- A: Many future networking applications will have a form of decentralized control, for scalability, fault-tolerance, and security.

- Example: P2P Spam-Filtering (Spamato-P2P).