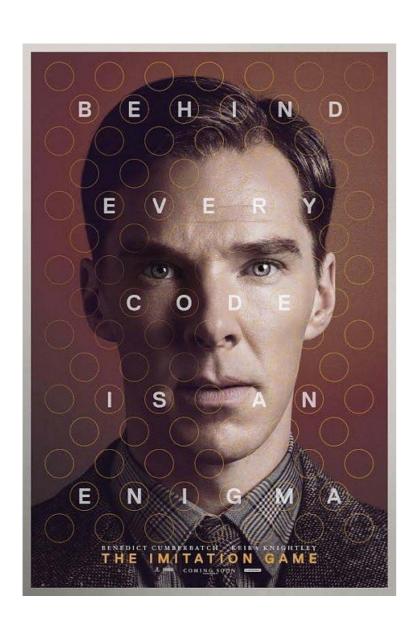
Automata & languages

A primer on the Theory of Computation



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Part 4 out of 5

Last week, we showed the equivalence of DFA, NFA and REX

We also started to look at non-regular languages

Pumping lemma

If A is a regular language, then there exist a number p s.t.

Any string in A whose length is at least p can be divided into three pieces xyz s.t.

- $xy^iz \in A$, for each i≥0 and
- |y| > 0 and
- $|xy| \leq p$

To prove that a language *A* is not regular:

- Assume that A is regular
- Since A is regular, it must have a pumping length p
- Find one string *s* in *A* whose length is at least *p*
- For any split *s*=*xyz*,

 Show that you cannot satisfy all three conditions
- 5 Conclude that *s* cannot be pumped

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- 1 Assume that A is regular
- 2 Since *A* is regular, it must have a pumping length *p*
- Find one string *s* in *A* whose length is at least *p*
- For any split *s=xyz*,
 Show that you cannot satisfy all three conditions
- 5 Conclude that s cannot be pumped \longrightarrow A is not regular

Wait...

What happens if A is a finite language?!

Pumping lemma

If A is a regular language, then there exist a number p s.t.

Any string in A whose length is at least p can be divided into three pieces xyz s.t.

- $xy^iz \in A$, for each i≥0 and
- |y| > 0 and
- $|xy| \leq p$

Pumping lemma

If *A* is a regular language, then there exist a number *p* s.t.

As we saw two weeks ago, all finite languages are regular...

So what's p?

p := len(longest_string) + 1
makes the lemma hold vacuously

Non-regular languages are not closed under most operations

if L_1 and L_2 are regular, then so are

L₁ U L₂
L₁ . L₂
L₁*

if L_1 and L_2 are not regular, then

$$\begin{array}{c|c} L_1 \cup L_2 \\ L_1 \cdot L_2 \\ L_1^* \end{array} \quad \begin{array}{c|c} \text{may or may not be} \\ \text{regular!} \\ \end{array}$$

 $(L_1)^C$ is not regular non RL are closed under complement

This week is all about

Context-Free Languages

a superset of Regular Languages