Automata & languages
A primer on the Theory of Computation

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October, 5 2017
Part 3 out of 5
Last week, we learned about closure and equivalence of regular languages
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The class of regular languages is closed under the

- union
- concatenation
- star

regular operations
The class of regular languages is closed under the regular operations if \( L_1 \) and \( L_2 \) are regular, then so are

- union \( L_1 \cup L_2 \)
- concatenation \( L_1 . L_2 \)
- star \( L_1^* \)
Last week, we learned about closure and equivalence of regular languages.
DFA $\cong$ NFA
This week we’ll look at REX, the third way of representing regular languages

\[ \text{DFA} \equiv \text{NFA} \]

REX
Are REX, NFA and DFA all equivalent?
We’ll then start asking ourselves whether all languages are regular

\[ L_1 \quad \{0^n1^n \mid n \geq 0\} \]

\[ L_2 \quad \{w \mid w \text{ has an equal number of 0s and 1s}\} \]

\[ L_3 \quad \{w \mid w \text{ has an equal number of occurrences of 01 and 10}\} \]

(only one of them actually is)
Advanced Automata
Thu Oct 5

1  Equivalence (the end)
   - DFA
   - NFA
   - Regular Expression

2  Non-regular languages

3  Context-free languages
Three tough languages

1) \( L_1 = \{0^n1^n \mid n \geq 0\} \)

2) \( L_2 = \{w \mid w \text{ has an equal number of 0s and 1s}\} \)

3) \( L_3 = \{w \mid w \text{ has an equal number of occurrences of } 01 \text{ and } 10 \text{ as substrings}\} \)
Three tough languages

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• In order to fully understand regular languages, we also must understand their limitations!
Pigeonhole principle

- Consider language $L$, which contains word $w \in L$.
- Consider an FA which accepts $L$, with $n < |w|$ states.
- Then, when accepting $w$, the FA must visit at least one state twice.
Pigeonhole principle

• Consider language L, which contains word \( w \in L \).
• Consider an FA which accepts L, with \( n < |w| \) states.
• Then, when accepting w, the FA must visit at least one state twice.

• This is according to the pigeonhole (a.k.a. Dirichlet) principle:
  – If \( m > n \) pigeons are put into \( n \) pigeonholes, there's a hole with more than one pigeon.
  – That’s a pretty fancy name for a boring observation...
Languages with unbounded strings

• Consequently, regular languages with unbounded strings can only be recognized by FA (finite! bounded!) automata if these long strings loop.

• The FA can enter the loop once, twice, ..., and not at all.
• That is, language L contains all \{xz, xyz, xy^2z, xy^3z, \ldots \}. 

Pumping Lemma

- Theorem:

  Given a regular language $L$, there is a number $p$ (the pumping number) such that:
  any string $u$ in $L$ of length $\geq p$ is pumpable within its first $p$ letters.
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any string $u$ in $L$ of length $\geq p$ is pumpable within its first $p$ letters.

• A string $u \in L$ with $|u| \geq p$ is pumpable if it can be split in 3 parts $xyz$ s.t.:
  – $|y| \geq 1$ (mid-portion $y$ is non-empty)
  – $|xy| \leq p$ (pumping occurs in first $p$ letters)
  – $xy^iz \in L$ for all $i \geq 0$ (can pump $y$-portion)
Pumping Lemma

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• If there is no such $p$, then the language is not regular
Pumping Lemma Example

• Let \( L \) be the language \( \{0^n1^n \mid n \geq 0\} \)

• Assume (for the sake of contradiction) that \( L \) is regular

• Let \( p \) be the pumping length. Let \( u \) be the string \( 0^p1^p \).

• Let’s check string \( u \) against the pumping lemma:

• “In other words, for all \( u \in L \) with \( |u| \geq p \) we can write:
  – \( u = xyz \) (\( x \) is a prefix, \( z \) is a suffix)
  – \( |y| \geq 1 \) (mid-portion \( y \) is non-empty)
  – \( |xy| \leq p \) (pumping occurs in first \( p \) letters)
  – \( xy^iz \in L \) for all \( i \geq 0 \) (can pump \( y \)-portion)”
Let’s make the example a bit harder...

- Let $L$ be the language $\{w \mid w$ has an equal number of 0s and 1$s\}$

- Assume (for the sake of contradiction) that $L$ is regular
- Let $p$ be the pumping length. Let $u$ be the string $0^p1^p$.
- Let’s check string $u$ against the pumping lemma:

  “In other words, for all $u \in L$ with $|u| \geq p$ we can write:
  
  - $u = xyz$ (x is a prefix, z is a suffix)
  - $|y| \geq 1$ (mid-portion $y$ is non-empty)
  - $|xy| \leq p$ (pumping occurs in first $p$ letters)
  - $xy^iz \in L$ for all $i \geq 0$ (can pump $y$-portion)”
Now you try...

• Is $L_1 = \{ww \mid w \in (0 \cup 1)^*\}$ regular?

• Is $L_2 = \{1^n \mid n \text{ being a prime number }\}$ regular?
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Part 1  regular language
Part 2  context-free language
Part 3  turing machine
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regular language
context-free language
turing machine

Part 2
Motivation

• Why is a language such as \( \{0^n1^n \mid n \geq 0\} \) not regular?!?

• It’s really simple! All you need to keep track is the number of 0’s...

• In this chapter we first study context-free grammars
  – More powerful than regular languages
  – Recursive structure
  – Developed for human languages
  – Important for engineers (parsers, protocols, etc.)
Example

- Palindromes, for example, are not regular.
- But there is a pattern.
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- Q: If you have one palindrome, how can you generate another?
- A: Generate palindromes recursively as follows:
  - Base case: $\varepsilon$, 0 and 1 are palindromes.
  - Recursion: If $x$ is a palindrome, then so are $0x0$ and $1x1$. 
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- Notation: $x \rightarrow \varepsilon | 0 | 1 | 0x0 | 1x1$.
  - Each pipe ("|") is an or, just as in UNIX regexp’s.
  - In fact, all palindromes can be generated from $\varepsilon$ using these rules.
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• Notation: \( x \rightarrow \varepsilon \mid 0 \mid 1 \mid 0x0 \mid 1x1 \).
  – Each pipe (“\(|\)”) is an or, just as in UNIX regexp’s.
  – In fact, all palindromes can be generated from \( \varepsilon \) using these rules.

• Q: How would you generate 11011011?
Context Free Grammars (CFG): Definition

- **Definition:** A context free grammar consists of \((V, \Sigma, R, S)\) with:
  - \(V\): a finite set of **variables** (or symbols, or non-terminals)
  - \(\Sigma\): a finite set of **terminals** (or the alphabet)
  - \(R\): a finite set of **rules** (or productions)
    - of the form \(v \rightarrow w\) with \(v \in V\), and \(w \in (\Sigma \cup V)^*\)
    - (read: “\(v\) yields \(w\)” or “\(v\) produces \(w\)”)
  - \(S \in V\): the **start symbol**.
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    of the form \(v \rightarrow w\) with \(v \in V\), and \(w \in (\Sigma_e \cup V)^*\)
    (read: “\(v\) yields \(w\)” or “\(v\) produces \(w\)”)
  – \(S \in V\): the start symbol.

• Q: What are \((V, \Sigma, R, S)\) for our palindrome example?
Derivations and Language

• Definition: The derivation symbol “⇒” (read “1-step derives” or “1-step produces”) is a relation between strings in $(\Sigma \cup \mathcal{V})^*$. We write $x \Rightarrow y$ if $x$ and $y$ can be broken up as $x = svt$ and $y = swt$ with $v \rightarrow w$ being a production in $R$. 
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- **Definition:** The derivation symbol “⇒*”, (read “derives” or “produces” or “yields”) is a relation between strings in \((\Sigma \cup V)^*\). We write \(x \Rightarrow^* y\) if there is a sequence of 1-step productions from \(x\) to \(y\). I.e., there are strings \(x_i\) with \(i\) ranging from 0 to \(n\) such that \(x = x_0, y = x_n\) and \(x_0 \Rightarrow x_1, x_1 \Rightarrow x_2, x_2 \Rightarrow x_3, \ldots, x_{n-1} \Rightarrow x_n\).
Derivations and Language

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• Definition: Let \(G\) be a context-free grammar. The context-free language (CFL) generated by \(G\) is the set of all terminal strings which are derivable from the start symbol. Symbolically: \(L(G) = \{w \in \Sigma^* | S \Rightarrow^* w\}\)
Example: Infix Expressions

- Infix expressions involving \{+, \times, a, b, c, (, )\}
- \(E\) stands for an expression (most general)
- \(F\) stands for factor (a multiplicative part)
- \(T\) stands for term (a product of factors)
- \(V\) stands for a variable: \(a, b,\) or \(c\)

- Grammar is given by:
  - \(E \rightarrow T \mid E + T\)
  - \(T \rightarrow F \mid T \times F\)
  - \(F \rightarrow V \mid (E)\)
  - \(V \rightarrow a \mid b \mid c\)

- Convention: Start variable is the first one in grammar \(E\)
Example: Infix Expressions

• Consider the string $u$ given by $a \times b + (c + (a + c))$
• This is a valid infix expression. Can be generated from $E$.

1. A sum of two expressions, so first production must be $E \Rightarrow E + T$
2. Sub-expression $a \times b$ is a product, so a term so generated by sequence $E + T \Rightarrow T + T \Rightarrow T \times F + T \Rightarrow^* a \times b + T$
3. Second sub-expression is a factor only because a parenthesized sum. $a \times b + T \Rightarrow a \times b + F \Rightarrow a \times b + (E) \Rightarrow a \times b + (E + T) \ldots$
4. $E \Rightarrow E + T \Rightarrow T + T \Rightarrow T \times F + T \Rightarrow F \times F + T \Rightarrow V \times F + T \Rightarrow a \times F + T \Rightarrow a \times V + T \Rightarrow a \times b + T \Rightarrow a \times b + F \Rightarrow a \times b + (E) \Rightarrow a \times b + (E + T) \Rightarrow a \times b + (T + T) \Rightarrow a \times b + (F + T) \Rightarrow a \times b + (V + T) \Rightarrow a \times b + (c + T) \Rightarrow a \times b + (c + F) \Rightarrow a \times b + (c + (E)) \Rightarrow a \times b + (c + (E + T)) \Rightarrow a \times b + (c + (T + T)) \Rightarrow a \times b + (c + (F + T)) \Rightarrow a \times b + (c + (a + T)) \Rightarrow a \times b + (c + (a + F)) \Rightarrow a \times b + (c + (a + V)) \Rightarrow a \times b + (c + (a + c))$
Left- and Right-most derivation

- The derivation on the previous slide was a so-called **left-most derivation**.

- In a **right-most derivation**, the variable most to the right is replaced.
  \[-E \Rightarrow E+T \Rightarrow E+F \Rightarrow E+(E) \Rightarrow E+(E+T) \Rightarrow etc.\]
Ambiguity

• There can be a lot of ambiguity involved in how a string is derived.

• Another way to describe a derivation in a unique way is using derivation trees.
Derivation Trees

- In a **derivation tree** (or parse tree) each node is a symbol. Each parent is a variable whose children spell out the production from left to right. For example, $v \rightarrow abcdefg$:

```
  v
 /\```
```
a b c d e f g```

- The root is the start variable.
- The leaves spell out the derived string from left to right.
Derivation Trees

- On the right, we see a derivation tree for our string \(a \times b + (c + (a + c))\)

- Derivation trees help understanding semantics! You can tell how expression should be evaluated from the tree.
Ambiguity

\[
\begin{align*}
\langle \text{sentence} \rangle & \rightarrow \langle \text{action} \rangle \mid \langle \text{action} \rangle \text{ with } \langle \text{subject} \rangle \\
\langle \text{action} \rangle & \rightarrow \langle \text{subject} \rangle \langle \text{activity} \rangle \\
\langle \text{subject} \rangle & \rightarrow \langle \text{noun} \rangle \mid \langle \text{noun} \rangle \text{ and } \langle \text{subject} \rangle \\
\langle \text{activity} \rangle & \rightarrow \langle \text{verb} \rangle \mid \langle \text{verb} \rangle \langle \text{object} \rangle \\
\langle \text{noun} \rangle & \rightarrow \text{Hannibal} \mid \text{Clarice} \mid \text{rice} \mid \text{onions} \\
\langle \text{verb} \rangle & \rightarrow \text{ate} \mid \text{played} \\
\langle \text{prep} \rangle & \rightarrow \text{with} \mid \text{and} \mid \text{or} \\
\langle \text{object} \rangle & \rightarrow \langle \text{noun} \rangle \mid \langle \text{noun} \rangle \langle \text{prep} \rangle \langle \text{object} \rangle
\end{align*}
\]

- Clarice played with Hannibal
- Clarice ate rice with onions
- Hannibal ate rice with Clarice

Q: Are there any suspect sentences?
Ambiguity

• A: Consider “Hannibal ate rice with Clarice”

• This could either mean
  – Hannibal and Clarice ate rice together.
  – Hannibal ate rice and ate Clarice.

• This ambiguity arises from the fact that the sentence has two different parse-trees, and therefore two different interpretations:
Hannibal and Clarice Ate

sentence

action

subject

noun

verb

object

Clarice

Hannibal

ate

with
Hannibal the Cannibal
Ambiguity: Definition

• Definition:

A string $x$ is said to be ambiguous relative the grammar $G$ if there are two essentially different ways to derive $x$ in $G$.
  – $x$ admits two (or more) different parse-trees
  – equivalently, $x$ admits different left-most [resp. right-most] derivations.

• A grammar $G$ is said to be ambiguous if there is some string $x$ in $L(G)$ which is ambiguous.
Ambiguity: Definition

• Definition:

A string $x$ is said to be **ambiguous** relative the grammar $G$ if there are two essentially different ways to derive $x$ in $G$.
  – $x$ admits two (or more) different parse-trees
  – equivalently, $x$ admits different left-most [resp. right-most] derivations.

• A grammar $G$ is said to be **ambiguous** if there is some string $x$ in $L(G)$ which is ambiguous.

• **Question**: Is the grammar $S \rightarrow ab \mid ba \mid aSb \mid bSa \mid SS$ ambiguous?
  – What language is generated?
Ambiguity

- Answer: $L(G) =$ the language with equal no. of $a'$ s and $b'$ s
- Yes, the language is ambiguous:
CFG’s: Proving Correctness

- The recursive nature of CFG’s means that they are especially amenable to correctness proofs.

- For example let’s consider the grammar
  \[ G = (S \rightarrow \varepsilon \mid ab \mid ba \mid aSb \mid bSa \mid SS) \]

- We claim that \( L(G) = L = \{ x \in \{a,b\}^* \mid n_a(x) = n_b(x) \} \),
  where \( n_a(x) \) is the number of \( a' \)s in \( x \), and \( n_b(x) \) is the number of \( b' \)s.

- **Proof**: To prove that \( L = L(G) \) is to show both inclusions:
  
  1. \( L \subseteq L(G) \): Every string in \( L \) can be generated by \( G \).
  2. \( L \supseteq L(G) \): \( G \) only generate strings of \( L \).
     - This part is easy, so we concentrate on part i.
Proving $L \subseteq L(G)$

- $L \subseteq L(G)$: Show that every string $x$ with the same number of $a$’s as $b$’s is generated by $G$. Prove by induction on the length $n = |x|$.
- Base case: The empty string is derived by $S \rightarrow \varepsilon$.
- Inductive hypothesis: Assume $n > 0$. Let $u$ be the smallest non-empty prefix of $x$ which is also in $L$.
  - Either there is such a prefix with $|u| < |x|$, then $x = uv$ whereas $v \in L$ as well, and we can use $S \rightarrow SS$ and repeat the argument.
  - Or $x = u$. In this case notice that $u$ can’t start and end in the same letter. If it started and ended with $a$ then write $x = ava$. This means that $v$ must have 2 more $b$’s than $a$’s. So somewhere in $v$ the $b$’s of $x$ catch up to the $a$’s which means that there’s a smaller prefix in $L$, contradicting the definition of $u$ as the smallest prefix in $L$. Thus for some string $v$ in $L$ we have $x = avb$ OR $x = bva$. We can use either $S \rightarrow aSb$ OR $S \rightarrow bS$.
Designing Context-Free Grammars

• As for regular languages this is a creative process.

• However, if the grammar is the union of simpler grammars, you can design the simpler grammars (with starting symbols $S_1$, $S_2$, respectively) first, and then add a new starting symbol/production $S \rightarrow S_1 \mid S_2$.

• If the CFG happens to be regular as well, you can first design the FA, introduce a variable/production for each state of the FA, and then add a rule $x \rightarrow ay$ to the CFG if $\delta(x,a) = y$ is in the FA. If a state $x$ is accepting in FA then add $x \rightarrow \varepsilon$ to CFG. The start symbol of the CFG is of course equivalent to the start state in the FA.

• There are quite a few other tricks. Try yourself…