



# Computer Systems

## Assignment 12

### 1 Game Theory

#### Quiz

#### 1.1 Selling a Franc

Form groups of two to three people. Every member of the group is a bidder in an auction for one (imaginary) franc. The franc is allocated to the highest bidder (for his/her last bid). Bids must be a multiple of CHF 0.05. This auction has a crux. Every bidder has to pay the amount of money he/she bid (last bid) – it does not matter if he/she gets the franc. Play the game!

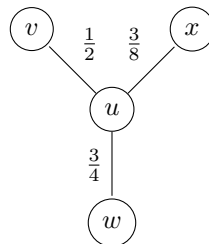
- a) Where did it all go wrong?
- b) What could the bidders have done differently?

#### Basic

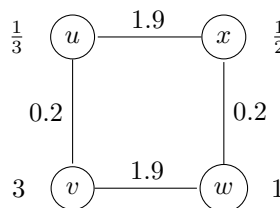
#### 1.2 Selfish Caching

For each of the following caching networks, compute the social optimum, the pure Nash equilibria, the price of anarchy (*PoA*) as well as the optimistic price of anarchy (*OPoA*):

- i.  $d_u = d_v = d_w = d_x = 1$



- ii. The demand is written next to a node.

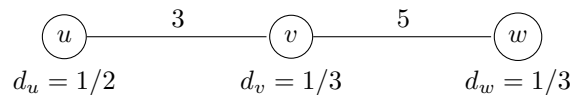


### 1.3 Selfish Caching with variable caching cost

The selfish caching model introduced in the lecture assumed that every peer incurs the same caching cost. However, this is a simplification of the reality. A peer with little storage space could experience a much higher caching cost than a peer who has terabytes of free disc space available. In this exercise, we omit the simplifying assumption and allow variable caching costs  $\alpha_i$  for node  $i$ .

What are the Nash Equilibria in the following caching networks given that

- i.  $\alpha_u = 1, \alpha_v = 2, \alpha_w = 2,$
- ii.  $\alpha_u = 3, \alpha_v = 3/2, \alpha_w = 3 ?$



Does any of the above instances have a dominant strategy profile? What is the PoA of each instance?

### Advanced

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#### 1.4 Matching Pennies

Tobias and Stephan like to gamble, and came up with the following game: Each of them secretly turns a penny to heads or tails. Then they reveal their choices simultaneously. If the pennies match Tobias gets both pennies, otherwise Stephan gets them.

Write down this 2-player game as a bi-matrix, and compute its (mixed) Nash equilibria!

#### 1.5 PoA Classes

The *PoA* of a class  $\mathcal{C}$  is defined as the maximum *PoA* over all instances in  $\mathcal{C}$ . Let

- $\mathcal{A}_{[a,b]}^n$  be the class of caching networks with  $n$  peers,  $a \leq \alpha_i \leq b$ ,  $d_i = 1$ , and each edge has weight 1,
- $\mathcal{W}_{[a,b]}^n$  be the class of networks with  $n$  peers,  $a \leq d_i \leq b$ ,  $\alpha_i = 1$ , and each edge has weight 1.

Show that  $PoA(\mathcal{A}_{[a,b]}^n) \leq \frac{b}{a} \cdot PoA(\mathcal{W}_{[\frac{1}{b}, \frac{1}{a}]}^n)$  for all  $n > 0$ .

## 2 Quorum Systems

### Quiz

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#### 2.1 The Resilience of a Quorum System

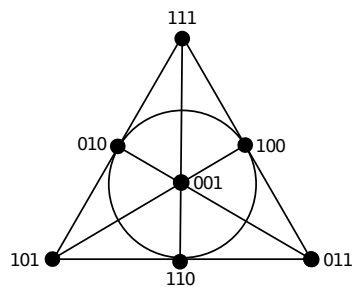
- Does a quorum system exist, which can tolerate that all nodes of a specific quorum fail? Give an example or prove its nonexistence.
- Consider the *nearly all* quorum system, which is made up of  $n$  different quorums, each containing  $n - 1$  servers. What is the resilience of this quorum system?
- Can you think of a quorum system that contains as many quorums as possible?  
*Note: the quorum system does not have to be minimal.*

### Basic

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#### 2.2 A Quorum System

Consider a quorum system with 7 nodes numbered from 001 to 111, in which each three nodes fulfilling  $x \oplus y = z$  constitute a quorum. In the following picture this quorum system is represented: All nodes on a line (such as 111, 010, 101) and the nodes on the circle (010, 100, 110) form a quorum.



- Of how many different quorums does this system consist of and what are its work and its load?
- Calculate its resilience  $f$ . Give an example where this quorum system does not work anymore with  $f + 1$  faulty nodes.

### Advanced

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#### 2.3 Uniform Quorum Systems

##### Definitions:

**s-Uniform:** A quorum system  $\mathcal{S}$  is *s-uniform* if every quorum in  $\mathcal{S}$  has exactly  $s$  elements.

**Balanced access strategy:** An access strategy  $Z$  for a quorum system  $\mathcal{S}$  is *balanced* if it satisfies  $L_Z(v_i) = L$  for all  $v_i \in V$ , for some value  $L$ .

**Claim:** An  $s$ -uniform quorum system  $\mathcal{S}$  reaches an optimal load with a balanced access strategy, if such a strategy exists.

- Describe in your own words why this claim is true.
- Prove the optimality of a balanced access strategy on an  $s$ -uniform quorum system.