



Ad Hoc And Sensor Networks

Exercise 7

Assigned: November 8, 2010

Due: November 15, 2010

1 Slotted Aloha

In this exercise we want to analyze ‘Slotted Aloha’ for the case that the number of stations n is not exactly known. We assume that in each time slot each station transmits with probability p . In the lecture you saw that the probability that the slot can be used (i.e. the probability that exactly one station transmits) is

$$\Pr(\text{success}) = n \cdot p(1 - p)^{n-1}.$$

If n is fixed, we can maximize the above expression and get the optimal p as shown in the lecture. Now assume that the only thing we know about n is $A \leq n \leq B$.

- Which value p maximizes $\Pr(\text{success})$ for the worst $n \in [A, B]$?
- What is this ‘worst case optimal’ value for p if $A = 100$ and $B = 200$?

2 Broadcast

Three students discuss the broadcasting problem with collision detection in graphs of constant diameter.

Student A claims that there is a deterministic protocol that allows to broadcast messages of length l in time $O(l)$. He says that it is possible since all nodes act synchronously and can detect collisions, which allows to transmit information one bit per round(slot) using the collision detection mechanism, i.e. detecting a transmission or collision in a slot means bit 1, detecting a free channel means 0.

Student B says that this is not possible because he can prove a lower bound of $\Omega(\log n)$ for deterministic algorithms, which can be much larger than the length of a message l in general. He says that this can be done in the same way as for the lower bound of n for the deterministic broadcast without collision detection for graphs of diameter 2, i.e. using golden and blue nodes in the middle layer.

Student C claims that A’s idea works in principle but all nodes need to know the length l of the message.

Who is right?

- If you believe A is right, give an algorithm that performs the broadcast.
- If you believe B is right, give a proof.
- If you believe C is right, describe an algorithm given that all nodes know the message length l and explain why the message length l is needed.