



# Ad Hoc And Sensor Networks

## Exercise 7

Assigned: November 8, 2010

Due: November 15, 2010

### 1 Slotted Aloha

In this exercise we want to analyze ‘Slotted Aloha’ for the case that the number of stations  $n$  is not exactly known. We assume that in each time slot each station transmits with probability  $p$ . In the lecture you saw that the probability that the slot can be used (i.e. the probability that exactly one station transmits) is

$$\Pr(\text{success}) = n \cdot p(1 - p)^{n-1}.$$

If  $n$  is fixed, we can maximize the above expression and get the optimal  $p$  as shown in the lecture. Now assume that the only thing we know about  $n$  is  $A \leq n \leq B$ .

- Which value  $p$  maximizes  $\Pr(\text{success})$  for the worst  $n \in [A, B]$ ?
- What is this ‘worst case optimal’ value for  $p$  if  $A = 100$  and  $B = 200$ ?

### 2 Broadcast

Three students discuss the broadcasting problem with collision detection in graphs of constant diameter.

Student A claims that there is a deterministic protocol that allows to broadcast messages of length  $l$  in time  $O(l)$ . He says that it is possible since all nodes act synchronously and can detect collisions, which allows to transmit information one bit per round(slot) using the collision detection mechanism, i.e. detecting a transmission or collision in a slot means bit 1, detecting a free channel means 0.

Student B says that this is not possible because he can prove a lower bound of  $\Omega(\log n)$  for deterministic algorithms, which can be much larger than the length of a message  $l$  in general. He says that this can be done in the same way as for the lower bound of  $n$  for the deterministic broadcast without collision detection for graphs of diameter 2, i.e. using golden and blue nodes in the middle layer.

Student C claims that A’s idea works in principle but all nodes need to know the length  $l$  of the message.

Who is right?

- If you believe A is right, give an algorithm that performs the broadcast.
- If you believe B is right, give a proof.
- If you believe C is right, describe an algorithm given that all nodes know the message length  $l$  and explain why the message length  $l$  is needed.