

Discrete Event Systems

Exercise session #4



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1. Context-Free Grammars

Give context-free grammars for the following languages over the alphabet $\Sigma = \{0, 1\}$:

- A) $L_1 = \{w \mid \text{the length of } w \text{ is odd}\}$
- B) $L_2 = \{w \mid \text{contains more 1s than 0s}\}$

1. Context-Free Grammars

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$$V = \{X, A\},$$

$$\Sigma = \{0, 1\},$$

$$R =$$

$$S = X$$

1. Context-Free Grammars

A) $L_1 = \{w \mid \text{the length of } w \text{ is odd}\}$

$$V = \{X, A\},$$

$$\Sigma = \{0, 1\},$$

$$R = \left\{ \begin{array}{l} X \rightarrow XAX \mid A, \\ A \rightarrow 0 \mid 1 \end{array} \right\},$$

$$S = X$$

1. Context-Free Grammars

B) $L_2 = \{w \mid \text{contains more 1s than 0s}\}$

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$$S \rightarrow A1A$$

$$A \rightarrow AA \mid 1A0 \mid 0A1 \mid 1 \mid \varepsilon$$

1. Context-Free Grammars

B) $L_2 = \{w \mid \text{contains more 1s than 0s}\}$

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2. Regular and Context-Free Languages

- a) Consider the context-free grammar G with the production $S \rightarrow SS \mid 1S2 \mid \epsilon$. Describe the language $L(G)$ in words, and prove that $L(G)$ is not regular.

2. Regular and Context-Free Languages

- a) Consider the context-free grammar G with the production $S \rightarrow SS \mid 1S2 \mid \epsilon$. Describe the language $L(G)$ in words, and prove that $L(G)$ is not regular.

Assume $L(G)$ is regular.

We take $w = 1^p 0 2^p \in L(G)$,

$w = xyz$ with $|xy| \leq p$ and $|y| \geq 1$, because of $|xy| \leq p$, **xy can only consist of 1s**

According to the pumping lemma, we should have $x\dot{y}z \in L$

However, by choosing $i=0$ we delete at least one 1 and get a word $w' = 1^{p-|y|} 0 2^p$ with $|y| \geq 1$.

w' is not in $L(G)$ since it has fewer 1s than 2s.

This means that w is not pumpable and hence, $L(G)$ is not regular.

2. Regular and Context-Free Languages

- b) The regular languages are a subset of the context-free languages. Give the context-free grammar for an arbitrary language L that is regular.

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- b) The regular languages are a subset of the context-free languages. Give the context-free grammar for an arbitrary language L that is regular.

Since every regular language is also context-free, we can choose an arbitrary regular language.

$L = \{0^n1, n \geq 1\}$ is clearly regular.

A context-free grammar for this language uses only the production $S \rightarrow 0S \mid 1$.

3. Pumping Lemma Revisited

A) Determine whether the language $L = \{1^{n^2} \mid n \in \mathbb{N}\}$ is regular.
Prove your claim!

3. Pumping Lemma Revisited

- B) Consider a regular language L and a pumping number p such that every word $u \in L$ can be written as $u = xyz$ with $|xy| \leq p$ and $|y| \geq 1$ such that $xy^iz \in L$ for all $i \geq 0$.

Can you use the pumping number p to determine the number of states of a minimal DFA accepting L ? What about the number of states of the corresponding NFA?

3. Pumping Lemma Revisited

$$\Sigma = \{a_1, a_2, \dots, a_n\}$$

$$L = \bigcup_{i=1}^n a_i^* = a_1^* \cup a_2^* \cup \dots \cup a_n^*$$

Minimum Pumping length 1,
but minimum DFA $n+2$