Principles of Distributed Computing
Exercise 9

1 Communication Complexity of Set Disjointness

In the lecture we studied the communication complexity of the equality function. Now we consider the disjointness function: Alice and Bob are given subsets $X, Y \subseteq \{1, \ldots, k\}$ and need to determine whether they are disjoint. Each subset can be represented by a string. E.g., we define the $i$th bit of $x \in \{0, 1\}^k$ as $x_i := 1$ if $i \in X$ and $x_i := 0$ if $i \notin X$. Now define disjointness of $X$ and $Y$ as:

$$\text{DISJ}(x, y) := \begin{cases} 0 : \text{there is an index } i \text{ such that } x_i = y_i = 1 \\ 1 : \text{else} \end{cases}$$

a) Write down $M^{\text{DISJ}}$ for the DISJ-function when $k = 3$.

b) Use the matrix obtained in a) to provide a fooling set of size 4 for DISJ in case $k = 3$.

c) In general, prove that $CC(\text{DISJ}) = \Omega(k)$.

2 Distinguishing Diameter 2 from 4

In the lecture we stated that when the bandwidth of an edge is limited to $O(\log n)$, the diameter of a graph can be computed in $O(n)$. In this problem, we show that we can do faster in case we know that all networks/graphs on which we execute an algorithm have either diameter 2 or diameter 4. We start by partitioning the nodes into sets: Let $s := s(n)$ be a threshold and define the set of high degree nodes $H := \{v \in V \mid d(v) \geq s\}$ and the set of low degree nodes $L := \{v \in V \mid d(v) < s\}$. Next, we define: An $H$-dominating set $\text{DOM}$ is a subset $\text{DOM} \subseteq V$ of the nodes such that each node in $H$ is either in the set $\text{DOM}$ or adjacent to a node in the set $\text{DOM}$.

Note: We define $N_1(v)$ as the closed neighborhood of vertex $v$ ($v$ and its adjacent nodes).

Assume in the following, that we can compute an $H$-dominating set $\text{DOM}$ of size $\frac{n \log n}{s}$ in time $O(D)$.

a) What is the distributed runtime of Algorithm 2-vs-4 (stated next page)? In case you believe that the distributed implementation of a step is not known from the lecture, find a distributed implementation for this step! Hint: The runtime depends on $s$ and $n$.

b) Find a function $s := s(n)$ such that the runtime is minimized (in terms of $n$).

c) Prove that if the diameter is 2, then Algorithm 2-vs-4 always returns 2.

Now assume that the diameter of the network is 4 and that we know vertices $u$ and $v$ with distance 4 to each other.
Algorithm 1 “2-vs-4”. Input: $G$ with diameter 2 or 4 Output: diameter of $G$

1. if $L \neq \emptyset$ then
2. \hspace{0.5cm} choose $v \in L$
3. \hspace{0.5cm} compute a BFS tree from each vertex in $N_1(v)$
4. else
5. \hspace{0.5cm} compute an $H$-dominating set $\text{DOM}$
6. \hspace{0.5cm} compute a BFS tree from each vertex in $\text{DOM}$
7. end if
8. if all BFS trees have depth 2 or 1 then
9. \hspace{0.5cm} return 2
10. else
11. \hspace{0.5cm} return 4
12. end if

\textbf{d)} Prove that if the algorithm performs a BFS from at least one node $w \in N_1(u)$ it decides “the diameter is 4”.

\textbf{e)} In case $L \neq \emptyset$: Prove that the algorithm performs a BFS of depth at least 3 from some node $w$. \textbf{Hint: use d)}

\textbf{f)} In case $L = \emptyset$: Prove that the algorithm performs a BFS of depth at least 3 from some node $w$.

\textbf{g)} Give a high level idea, why you think that this does not violate the lower bound of $\Omega(n/\log n)$ presented in the lecture!

\textbf{h)} Assume $s = \frac{n}{2}$. Prove or disprove: If the diameter is 2, then Algorithm 2-vs-4 will always compute some BFS tree of depth exactly 2.