



Principles of Distributed Computing

Exercise 12: Sample Solution

We provide a solution for Task 2, which includes a solution for Task 1 as a special case. The construction is the following

- 1) Given G , compute a rooted spanning tree T .
- 2) Set $\alpha = (\epsilon n)/2k$.
- 3) For a node v , let $C(v)$ be the children of v in T .
- 4) The algorithm processes nodes from the bottom on, and computes a weight $W(v)$ for each node as follows.

$$W'(v) = 1 + \sum_{u \in C(v)} W(u).$$

- 5) If $W'(v) \geq \alpha$, then add v to detector set; and put $W(v) = 0$. Otherwise, let $W(v) = W'(v)$.

Now we prove that the detector set D output by this algorithm is a weak (ϵ, k) detection set.

Let \mathcal{C} be the set of connected components in $T \setminus D$. First, we note that every connected component in \mathcal{C} has size at most α . If the weight were greater than α , look at the node in this component that is the ancestor of all other nodes in the tree-scan order. Now this node has weight greater than α , which would imply it would have been added to D . A contradiction!

We now show that D is a detection set. So, consider any subset $Z \subset V$ of size at most k . If $Z \cap D$ is non-empty, then we are done. So, assume that $Z \cap D = \emptyset$. Consider two disjoint sets C_1, C_2 , each of size at least $\epsilon \cdot n$, in $G \setminus Z$. We will show that at least one of them contains a vertex of D , which is the witness to the failure.

If D intersects both C_1 and C_2 , we are done. So, assume that C_1 has empty intersection with D . Because $|C_1| \geq \epsilon \cdot n$, there are at least $2k$ components A_1, A_2, \dots, A_{2k} of \mathcal{C} that intersect C_1 . In order to disconnect $(C_1 \cap A_i)$ from $G \setminus C_1$ without including any detectors, the adversary must delete at least one element (vertex or edge) from A_i . Hence, he must delete at least $2k > k$ vertices, which contradicts the definition of Z .